

CS41B MACHINE

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CS 52 – Fall 2015

Admin

- Assignment 3
 - due Monday at 11:59pm
 - one small error in 5b (fast division) that's been fixed
- Midterm next Thursday in-class (10/1)
 - Comprehensive ☺
 - Closed books, notes, computers, etc.
 - **Except**, may bring up to 2 pages of notes
 - Practice problems posted
 - Also some practice problems in the Intro SML reading
 - Midterm review sessions Tuesday and Wednesday (Q&A session for midterm)

Midterm topics

SML

recursion

math

Midterm topics

- Basic syntax
- SML built-in types
- Defining function
 - pattern matching
- Function type signatures
- Recursion!
- map
- exceptions
- Defining datatypes
- addition, subtraction, multiplication manually and on list digits
- Numbers in different bases
- Binary number representation (first part of today's lecture)
- *NOT CS41B material*

Binary number revisited

What number does 1001 represent in binary?

Depends!
Is it a signed number or unsigned?
If signed, what convention are we using?

Twos complement

For a number with n digits high order bit represents -2^{n-1}

unsigned

2^3	2^2	2^1	2^0

signed
(twos complement)

-2^3	2^2	2^1	2^0

Twos complement

What number is it?

unsigned	1	0	0	1	9
	2^3	2^2	2^1	2^0	
signed (twos complement)	1	0	0	1	-7
	-2^3	2^2	2^1	2^0	

Twos complement

What number is it?

unsigned	1	1	1	1	15
	2^3	2^2	2^1	2^0	
signed (twos complement)	1	1	1	1	-1
	-2^3	2^2	2^1	2^0	

Twos complement

What number is it?

unsigned	1	1	0	0	12
	2^3	2^2	2^1	2^0	
signed (twos complement)	1	1	0	0	-4
	-2^3	2^2	2^1	2^0	

Twos complement

How many numbers can we represent with each approach using 4 bits?

16 (2^4) numbers, 0000, 0001, ..., 1111
Doesn't matter the representation!

unsigned				
	2^3	2^2	2^1	2^0
signed (twos complement)				
	-2^3	2^2	2^1	2^0

Twos complement

How many numbers can we represent with each approach using 32 bits?

$2^{32} \approx 4$ billion numbers

unsigned				
	2^3	2^2	2^1	2^0
signed (twos complement)				
	-2^3	2^2	2^1	2^0

Twos complement

What is the range of numbers we can represent for each approach?

unsigned: 0, 1, ... 15
signed: -8, -7, ..., 7

unsigned				
	2^3	2^2	2^1	2^0
signed (twos complement)				
	-2^3	2^2	2^1	2^0

binary representation	unsigned	
0000	0	
0001	1	
0010	?	
0011		
0100		
0101		
0110		
0111		
1000		
1001		
1010		
1011		
1100		
1101		
1110		
1111		

binary representation	unsigned	twos complement
0000	0	?
0001	1	
0010	2	
0011	3	
0100	4	
0101	5	
0110	6	
0111	7	
1000	8	
1001	9	
1010	10	
1011	11	
1100	12	
1101	13	
1110	14	
1111	15	

binary representation	unsigned	twos complement
0000	0	0
0001	1	1
0010	2	2
0011	3	3
0100	4	4
0101	5	5
0110	6	6
0111	7	7
1000	8	?
1001	9	
1010	10	
1011	11	
1100	12	
1101	13	
1110	14	
1111	15	

binary representation	unsigned	twos complement
0000	0	0
0001	1	1
0010	2	2
0011	3	3
0100	4	4
0101	5	5
0110	6	6
0111	7	7
1000	8	-8
1001	9	?
1010	10	
1011	11	
1100	12	
1101	13	
1110	14	
1111	15	

binary representation	unsigned	twos complement
0000	0	0
0001	1	1
0010	2	2
0011	3	3
0100	4	4
0101	5	5
0110	6	6
0111	7	7
1000	8	-8
1001	9	-7
1010	10	-6
1011	11	-5
1100	12	-4
1101	13	-3
1110	14	-2
1111	15	-1

binary representation	unsigned	twos complement
0000	0	0
0001	1	1
0010	2	2
0011	3	3
0100	4	4
0101	5	5
0110	6	6
0111	7	7
1000	8	-8
1001	9	-7
1010	10	-6
1011	11	-5
1100	12	-4
1101	13	-3
1110	14	-2
1111	15	-1

How can you tell if a number is negative?

binary representation	unsigned	twos complement
0000	0	0
0001	1	1
0010	2	2
0011	3	3
0100	4	4
0101	5	5
0110	6	6
0111	7	7
1000	8	-8
1001	9	-7
1010	10	-6
1011	11	-5
1100	12	-4
1101	13	-3
1110	14	-2
1111	15	-1

High order bit!

A two's complement trick

You can also calculate the value of a negative number represented as twos complement as follows:

- Flip all of the bits (0 → 1 and 1 → 0)
- Add 1
- Resulting number is the magnitude of the original negative number



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- Resulting number is the magnitude of the original negative number

flip the bits
add 1

1110 → 0001 → 0010 → -2

Addition with two's complement numbers

$$\begin{array}{r} 0001 \\ + 0101 \\ \hline \end{array}$$

?

Addition with two's complement numbers

$$\begin{array}{r} \\ 0001 \\ + 0101 \\ \hline 0110 \end{array}$$

Addition with two's complement numbers

$$\begin{array}{r} 0110 \\ + 0101 \\ \hline \end{array}$$

?

Addition with twos complement numbers

$$\begin{array}{r} \overset{1}{0}110 \\ + 0101 \\ \hline 1011? \end{array}$$

Addition with twos complement numbers

$$\begin{array}{r} \overset{1}{0}110 \quad 6 \\ + 0101 \quad 5 \\ \hline 1011? \quad 11 \end{array}$$

Overflow! We cannot represent this number (it's too large)

Addition with twos complement numbers

$$\begin{array}{r} 0110 \\ + 1101 \\ \hline ? \end{array}$$

Addition with twos complement numbers

$$\begin{array}{r} \overset{1}{1} \overset{1}{0}110 \\ + 1101 \\ \hline 0011 \end{array}$$

Addition with two's complement numbers

ignore the last carry

$$\begin{array}{r}
 \begin{array}{c} \nearrow \\ 1 \end{array} \begin{array}{c} \nearrow \\ 1 \end{array} \\
 0110 \quad 6 \\
 + 1101 \quad -3 \\
 \hline
 0011 \quad 3
 \end{array}$$

Subtraction

Ideas?

Subtraction

- Negate the 2nd number (flip the bits and add 1)
- Add them!

A two's complement trick

You can also calculate the value of a negative number represented as two's complement as follows:

- Flip all of the bits (0 → 1 and 1 → 0)
- Add 1
- Resulting number is the magnitude of the original negative number

Why does this work?

A two's complement trick

You can also calculate the value of a negative number represented as two's complement as follows:

- ▣ Flip all of the bits (0 → 1 and 1 → 0)
- ▣ Add 1
- ▣ Resulting number is the magnitude of the original negative number

What is:

$$1101 + 0010 \quad ?$$

Hexadecimal numbers

Hexadecimal = base 16

What will be the digits?

Hexadecimal numbers

Hexadecimal = base 16

Digits

- 0
- 1
- 2
- ...
- 9
- a (10)
- b (11)
- c (12)
- d (13)
- e (14)
- f (15)

What number is 1ad?

Hexadecimal numbers

Hexadecimal = base 16

Digits

- 0
- 1
- 2
- ...
- 9
- a (10)
- b (11)
- c (12)
- d (13)
- e (14)
- f (15)

$$1 \quad a \quad d = 256 + 10 * 16 + 13 = 429$$

Hexadecimal numbers

Hexadecimal = base 16

Digits	
0	
1	Hexadecimal is common in CS.
2	Why?
...	
9	
a (10)	
b (11)	
c (12)	
d (13)	
e (14)	
f (15)	

Hexadecimal numbers

Hexadecimal = base 16

Digits	
0	
1	Hexadecimal is common in CS.
2	4 bits = 1 hexadecimal digit
...	
9	
a (10)	What is the following number in hex?
b (11)	
c (12)	
d (13)	
e (14)	
f (15)	

0110 1101 0101 0001

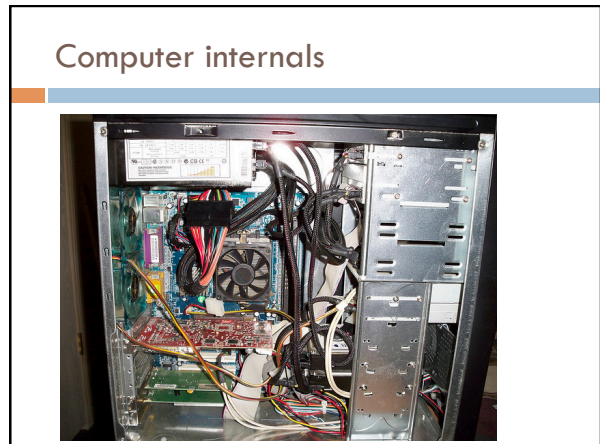
Hexadecimal numbers

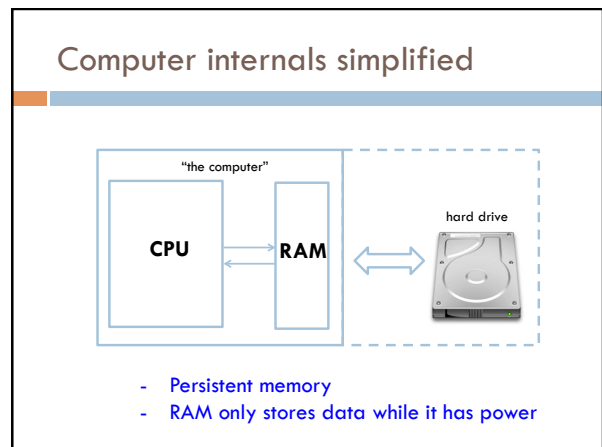
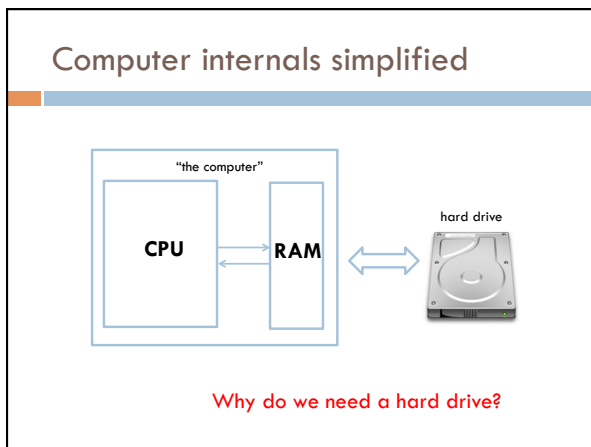
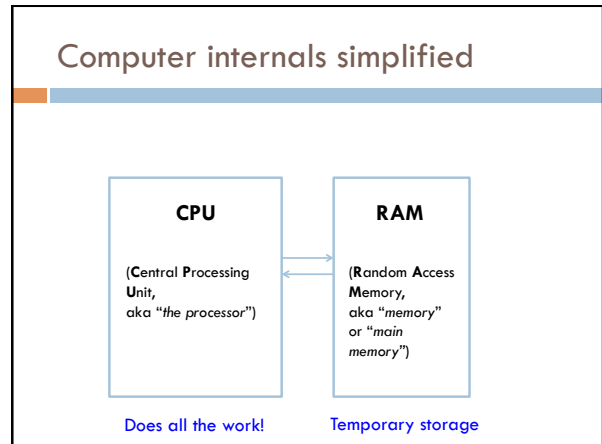
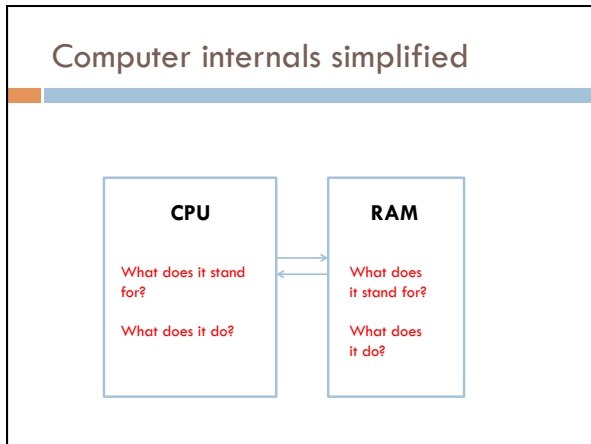
Hexadecimal = base 16

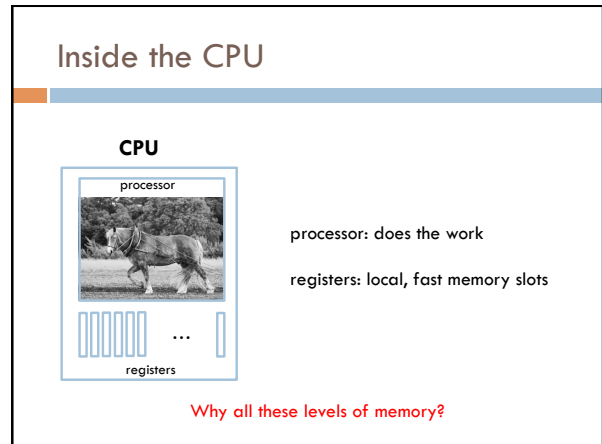
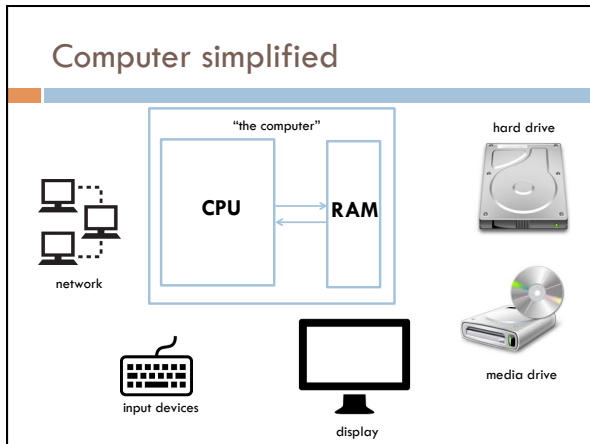
Digits	
0	
1	Hexadecimal is common in CS.
2	4 bits = 1 hexadecimal digit
...	
9	
a (10)	What is the following number in hex?
b (11)	
c (12)	
d (13)	
e (14)	
f (15)	

0110 1101 0101 0001

6 d 5 1

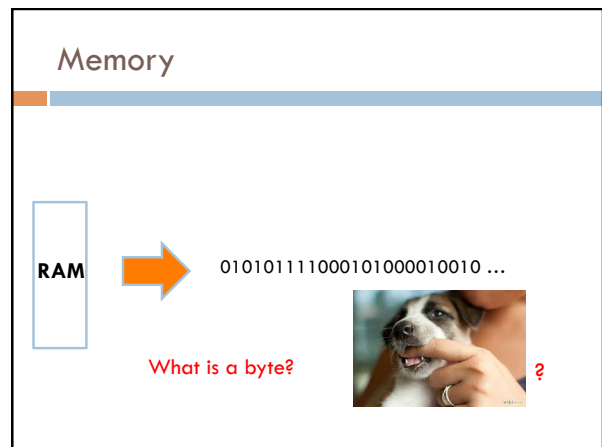






Memory speed

operation	access time	times slower than register access	for comparison ...
register	0.3 ns	1	1 s
RAM	120 ns	400	6 min
Hard disk	1ms	~million	1 month
google.com	0.4s	~billion	30 years



Memory

RAM → 01010111 10001010 00010010 ...

byte = 8 bits
byte is abbreviated as B

My laptop has 16GB (gigabytes) of memory. How many bits is that?

Memory sizes

	bits
byte	8
kilobyte (KB)	2^{10} bytes = ~8,000
megabyte (MB)	2^{20} = ~ 8 million
gigabyte (GB)	2^{30} = ~8 billion

My laptop has 16GB (gigabytes) of memory. How many bits is that?

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gigabyte (GB)	2^{30} = ~8 billion

~128 billion bits!

Memory

RAM →

address	
0	01010111
1	10001010
2	00010010
3	01011010
...	...

Memory is byte addressable

Memory

address	value
0	01010111
1	10001010
2	00010010
3	01011010
...	...

Memory is organized into "words", which is the most common functional unit

Memory

address	32-bit words
0	10101011 10001010 00010010 01011010
4	11001011 00001110 01010010 01010110
8	10111011 10010010 00000000 01110100
...	...

Most modern computers use 32-bit (4 byte) or 64-bit (8 byte) words

Memory in the CS41B Machine

address	16-bit words
0	10101011 10001010
2	00010010 01011010
4	11001011 00001110
...	...

We'll use 16-bit words for our model (the CS41B machine)

CS41B machine


CPU

- processor (includes a picture of a horse)
- registers (four vertical bars)
- ic: instruction counter (location in memory of the next instruction in memory)
- r0: holds the value 0 (read only)
- r1, r2, r3: general purpose read/write registers

CS41B instructions

CPU

processor



registers

What types of operations might we want to do (think really basic)?

abbreviation	arguments	action
Register Instructions		
mov	RR-	dest = src0
neg	RR-	dest = -src0
add	RRR	dest = src0 + src1
sub	RRR	dest = src0 - src1
adc	RRS	dest = src0 + arg
sbc	RRS	dest = src0 - arg

abbreviation	arguments	action
Register Instructions		
mov	RR-	dest = src0
neg	RR-	dest = -src0
add	RRR	dest = src0 + src1
sub	RRR	dest = src0 - src1
adc	RRS	dest = src0 + arg
sbc	RRS	dest = src0 - arg

operation name
(always three characters)

abbreviation	arguments	action
Register Instructions		
mov	RR-	dest = src0
neg	RR-	dest = -src0
add	RRR	dest = src0 + src1
sub	RRR	dest = src0 - src1
adc	RRS	dest = src0 + arg
sbc	RRS	dest = src0 - arg

operation arguments
R = register (e.g. r0)
S = signed number (byte)

abbreviation	arguments	action
Register Instructions		
mov	RR-	dest = src0
neg	RR-	dest = -src0
add	RRR	dest = src0 + src1
sub	RRR	dest = src0 - src1
adc	RRS	dest = src0 + arg
sbc	RRS	dest = src0 - arg

operation function
 dest = first register
 src0 = second register
 src1 = third register
 arg = number/argument

add r1 r2 r3

What does this do?

abbreviation	arguments	action
Register Instructions		
mov	RR-	dest = src0
neg	RR-	dest = -src0
add	RRR	dest = src0 + src1
sub	RRR	dest = src0 - src1
adc	RRS	dest = src0 + arg
sbc	RRS	dest = src0 - arg

add r1 r2 r3

$r1 = r2 + r3$

Add contents of registers r2 and r3 and store the result in r1

abbreviation	arguments	action
Register Instructions		
mov	RR-	dest = src0
neg	RR-	dest = -src0
add	RRR	dest = src0 + src1
sub	RRR	dest = src0 - src1
adc	RRS	dest = src0 + arg
sbc	RRS	dest = src0 - arg

adc r2 r1 10

What does this do?

abbreviation	arguments	action
Register Instructions		
mov	RR-	dest = src0
neg	RR-	dest = -src0
add	RRR	dest = src0 + src1
sub	RRR	dest = src0 - src1
adc	RRS	dest = src0 + arg
sbc	RRS	dest = src0 - arg

adc r2 r1 10

$r2 = r1 + 10$

Add 10 to the contents of register r1 and store in r2

abbreviation	arguments	action
--------------	-----------	--------

Register Instructions

mov	RR-	dest = src0
neg	RR-	dest = -src0
add	RRR	dest = src0 + src1
sub	RRR	dest = src0 - src1
adc	RRS	dest = src0 + arg
sbc	RRS	dest = src0 - arg

adc r1 r0 8

neg r2 r1

sub r2 r1 r2

What number is in r2?

abbreviation	arguments	action
--------------	-----------	--------

Register Instructions

mov	RR-	dest = src0
neg	RR-	dest = -src0
add	RRR	dest = src0 + src1
sub	RRR	dest = src0 - src1
adc	RRS	dest = src0 + arg
sbc	RRS	dest = src0 - arg

adc r1 r0 8

$r1 = 8$

neg r2 r1

$r2 = -8, r1 = 8$

sub r2 r1 r2

$r2 = 16$

abbreviation	arguments	action
--------------	-----------	--------

Register Instructions

mov	RR-	dest = src0
neg	RR-	dest = -src0
add	RRR	dest = src0 + src1
sub	RRR	dest = src0 - src1
adc	RRS	dest = src0 + arg
sbc	RRS	dest = src0 - arg